# Was the Year 2000 a Leap Year? Step-wise Narrowing Theories with Metagol

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2016



(2014)

[2013]

2012

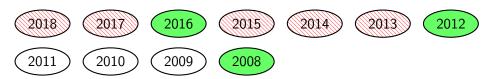


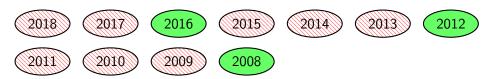


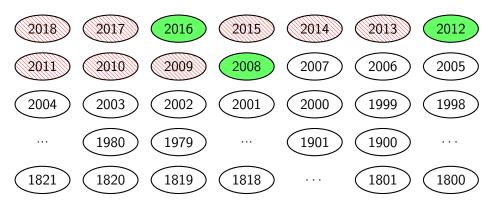
2014

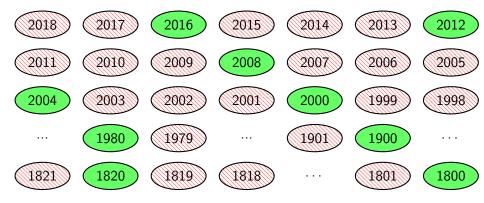
2013

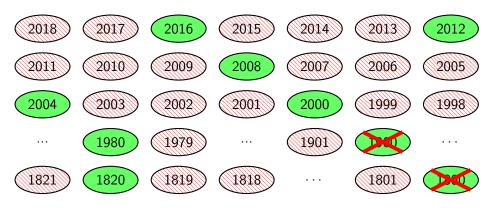
2012











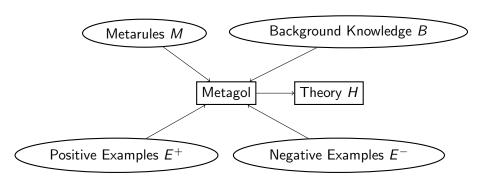
#### Definition [Richards, 2013 (p. 599)]

Every year that is exactly divisible by 4 is a leap year, *except* for years that are exactly divisible by 100, *but* these centurial years are leap years if they are exactly divisible by 400.

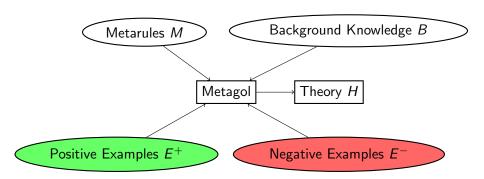
#### Logic Program

```
leapyear(X) \leftarrow divisible(X,4), not divisible(X,100). leapyear(X) \leftarrow divisible(X,400).
```

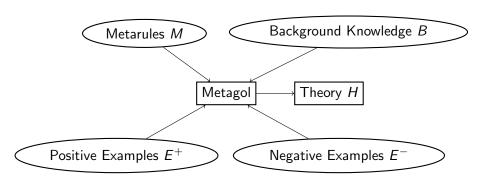
## Metagol



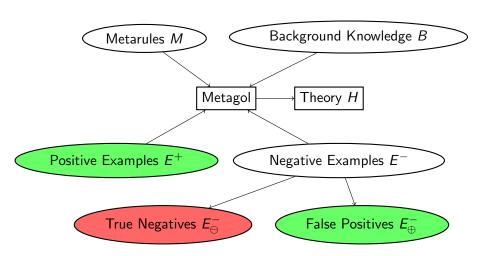
## Metagol



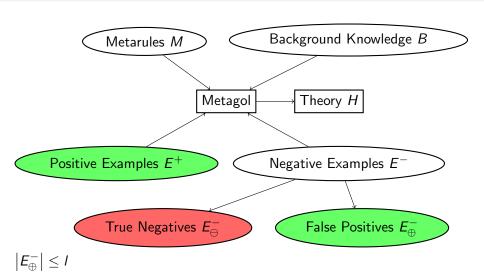
## Relaxed Metagol



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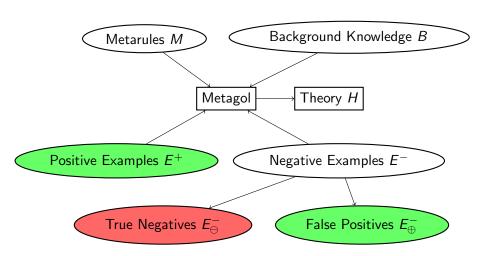
## Logic Program

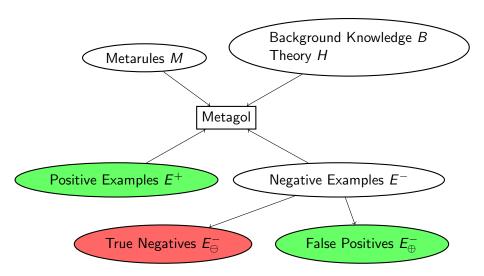
 $leapyear(X) \leftarrow divisible(X,4)$ .

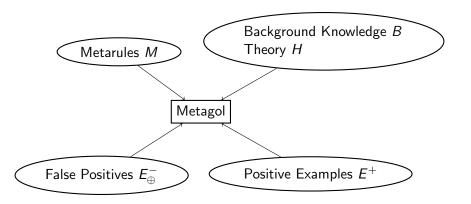
#### False Positives

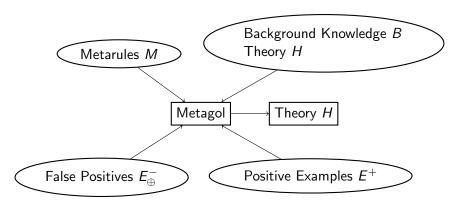
1900, 1800, 1700, 1500, 1400, ...

 $\Rightarrow$  99.25% accuracy









## Step-wise Narrowed Theory

#### Definition (Step-wise Narrowed Theory)

Every higher-order datalog program is a step-wise narrowed theory. Let H be a higher-order datalog program,  $\Phi$  a mapping between predicate symbols and S a step-wise narrowed theory. Then,  $\langle H, \Phi, S \rangle$  is also a step-wise narrowed theory.

## Step-wise Narrowed Theory $leapyear(A) \leftarrow divisible(A,4).,$ $leapyear \mapsto \_leapyear$ , $_{\text{leapyear}}(A) \leftarrow \text{divisible}(A,100).,$ $_{\text{leapyear}} \mapsto _{\text{leapyear}}$ $_$ leapyear(A) ← divisible(A,400).

```
\label{eq:leapyear} \begin{split} & | \mathsf{leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A},\!4)., \\ & | \mathsf{leapyear} \mapsto \_\mathsf{leapyear}, \\ & \langle \\ & \_\mathsf{leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A},\!100)., \\ & \_\mathsf{leapyear} \mapsto \_\_\mathsf{leapyear}, \\ & \_\mathsf{leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A},\!400). \\ & \rangle \\ \end{split}
```

```
\label{eq:leapyear} \begin{array}{l} \mathsf{leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A},\!\mathsf{4})., \\ \mathsf{leapyear} \mapsto \_\mathsf{leapyear}, \\ \langle \\ \_\mathsf{leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A},\!100), \ \mathsf{not}(\_\_\mathsf{leapyear}(\mathsf{A})). \\ \_\_\mathsf{leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A},\!400). \\ \rangle \\ \end{array}
```

```
\begin{split} & \mathsf{leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A},\!\mathsf{4}), \ \mathsf{not}(\mathsf{\_leapyear}(\mathsf{A})). \\ & \mathsf{\_leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A},\!\mathsf{100}), \ \mathsf{not}(\mathsf{\_\_leapyear}(\mathsf{A})). \\ & \mathsf{\_leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A},\!\mathsf{400}). \end{split}
```

#### Flattened Theory

```
\begin{split} & \mathsf{leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A}, \mathsf{4}), \ \mathsf{not}(\mathsf{\_leapyear}(\mathsf{A})). \\ & \mathsf{\_leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A}, \mathsf{100}), \ \mathsf{not}(\mathsf{\_\_leapyear}(\mathsf{A})). \\ & \mathsf{\_leapyear}(\mathsf{A}) \leftarrow \mathsf{divisible}(\mathsf{A}, \mathsf{400}). \end{split}
```

#### Caution

Only for non-recursive theories!

## Compare Metagol and Metagol<sub>SN</sub>

#### Experimental Setup

Domain leapyear

#### Training data

- years 1582-2018
- randomly samples 20%, 40%, ..., 100%

Test data years 2019-3018

Repititions 10x

Learning time out 30 min

## Compare Metagol and Metagol<sub>SN</sub>

#### Background Knowledge

divisible/2 where divisible(X, Y) holds if and only if the integer Y divides the integer X exactly and

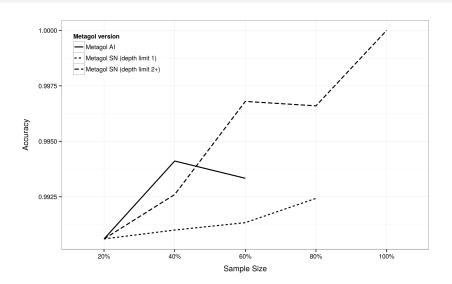
not\_divisible/2 where  $not_divisible(X, Y)$  holds if and only if X and Y are integers and divisible(X, Y) does not hold.

Y's could be choosen from all divisors of all examples.

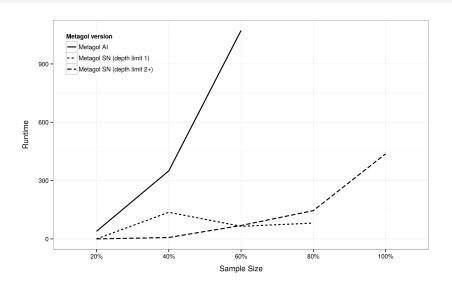
#### Metarules

Const 
$$P(A, B) \leftarrow$$
  
And  $P(A) \leftarrow Q(A), R(A)$   
Chain  $P(A, B) \leftarrow Q(A, C), R(C, B)$   
Curry  $P(A) \leftarrow Q(A, B)$ 

## Results



## Results



#### Conclusion

#### Summary

- Relaxed Metagol
- Allowed (some) negation in Metagol
- Metagol<sub>SN</sub> is faster
- Metagol<sub>SN</sub> might have better results

#### Further Research

- Evaluate Metagol<sub>SN</sub> on more domains
- Use relaxed Metagol for pre-tests
- Use relaxed Metagol for anytime-algorithm

## Try it!

#### Available on GitHub



https://github.com/michael-siebers/metagol/tree/ilp2018

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## Questions / Comments

## Metagol Code

```
metagol(Pos,Neg,Prog) :-
  prove_all(Pos,[],Prog),
  prove_none(Neg,Prog).
prove_all([],Prog,Prog).
prove_all([Atom|Atoms],Prog1,Prog2) :-
  prove_one(Atom, Prog1, Prog3),
  prove_all(Atoms, Prog3, Prog2).
prove_one(Atom,Prog,Prog) :- call(Atom).
prove_one(Atom,Prog1,Prog2) :-
  metarule(Name, MetaSub, (Atom :- Body)),
  store(sub(Name, MetaSub), Prog1, Prog3),
  prove_all(Body,Prog3,Prog2).
```

## Metagol Code (Cont'd)

```
metagol(Pos,Neg,Prog) :-
   prove_all(Pos,[],Prog),
   prove_none(Neg,Prog).

prove_none([],Prog).
prove_none([Atom|Atoms],Prog) :-
   not(prove_one(Atom,Prog,Prog)),
   prove_none(Atoms,Prog).
```

## Metagol Relaxed

```
metagol_relaxed(Pos,Neg,Prog,FalsePos) :-
  prove_all(Pos,[],Prog),
  prove_some(Neg,Prog,FalsePos).
prove_some([],Prog,[]).
prove_some([Atom|Atoms],Prog,[Atom|Proven]) :-
  prove_some(Atoms, Prog, Proven),
  prove_one(Atom, Prog, Prog).
prove_some([Atom|Atoms],Prog,Proven) :-
  prove_some(Atoms, Prog, Proven),
  not(prove_one(Atom,Prog,Prog)).
```

## Metagol<sub>SN</sub>

```
metagol_sn(Pos,Neg,SNT) :-
  let P be the predicate symbol used in Pos,
  metagol_relaxed(Pos,Neg,Prog1,FalsePos),
  if FalsePos=∏
    SNT = Prog1
  else
    let PPrime be P prefixed with '_',
    let PosNext be FalsePos with P renamed to PPrime,
    let NegNext be Pos with P renamed to PPrime,
    assert_prog(Prog1),
    metagol_sn(PosNext,NegNext,Prog2),
    SNT=snt(Prog1, PPrime, Prog2).
```